

Towards Higher Observational Type Theory

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- ▶ **Higher** Observational Type Theory (**HOTT**):

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$$\text{Id}_{\text{Type}} A B := A \simeq B$$

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- + Efficient
- Higher inductive types only work in the empty context;
computation rules for higher constructors are typal

Why is OTT \rightarrow HOTT difficult?

Computation in OTT:

$$\text{coe} : (A_0 A_1 : \text{Set}) \rightarrow \text{Id}_{\text{Set}} A_0 A_1 \rightarrow A_0 \rightarrow A_1$$

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 - ▶ if the relation is the graph of a function: $f a = b \rightarrow f (\alpha_A a) = \alpha_B b$

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$$\begin{array}{ccc} & a_1^P & \\ & \rightarrow & \\ a_1 & & a_1 \\ \uparrow & & \uparrow \\ a_2 & & a_2 \\ \uparrow & & \uparrow \\ a_0 & & a_0 \\ & a_0^P & \end{array}$$

$$\begin{array}{ccc} & a_1 & \\ & \uparrow & \\ a_2 & & \\ & a_0 & \end{array}$$

$$\begin{array}{ccc} & a_2 & \\ & \rightarrow & \\ a_0 & & a_1 \\ \uparrow & & \uparrow \\ a_0^P & & a_1^P \\ & a_2 & \\ & \rightarrow & \\ & a_0 & \end{array}$$

Theories of internal parametricity

- ▶ Bernardy–Moulin (2012, 2013): higher dimensional syntax
- ▶ Bernardy–Coquand–Moulin (2015): substructural interval
- ▶ Nuyts–Devriese (2018): simpler semantics
- ▶ Cavallo–Harper (2021): parametricity and univalence combined
- ▶ Altenkirch–Chamoun–Kaposi–Shulman (2024): first structural theory
- ▶ Mike Shulman (2024–2026): the proof assistant Narya
- ▶ Sarah Reboullet’s thesis (2025): analysis of a problem in BCM 2015
- ▶ Jem Lord (2025): internal axiomatic parametricity

Informal Narya

$$\frac{A : \text{Type}}{A^P : A \rightarrow A \rightarrow \text{Type}}$$

$$\frac{a : A}{a^P : A^P a a}$$

Informal Narya

$$\frac{A : \text{Type}}{A^P : A \rightarrow A \rightarrow \text{Type}}$$

$$\frac{a : A}{a^P : A^P a a}$$

$$\frac{R : A \rightarrow B \rightarrow \text{Type}}{\text{Gel } R : \text{Type}^P A B}$$

$$\frac{A_2 : \text{Type}^P A_0 A_1}{A_2 : A_0 \rightarrow A_1 \rightarrow \text{Type}}$$

$$\text{Gel } R a b \cong R a b$$

Informal Narya

$$\frac{A : \text{Type}}{A^P : A \rightarrow A \rightarrow \text{Type}} \qquad \frac{a : A}{a^P : A^P a a}$$

$$\frac{R : A \rightarrow B \rightarrow \text{Type}}{\text{Gel } R : \text{Type}^P A B} \qquad \frac{A_2 : \text{Type}^P A_0 A_1}{A_2 : A_0 \rightarrow A_1 \rightarrow \text{Type}} \qquad \text{Gel } R a b \cong R a b$$

$$(A \rightarrow B)^P f_0 f_1 = \{a_0 a_1 : A\} \rightarrow A^P a_0 a_1 \rightarrow B^P (f_0 a_0) (f_1 a_1)$$

Informal Narya

$$\frac{A : \text{Type}}{A^P : A \rightarrow A \rightarrow \text{Type}} \qquad \frac{a : A}{a^P : A^P a a}$$

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$$(A \rightarrow B)^P f_0 f_1 = \{a_0 a_1 : A\} \rightarrow A^P a_0 a_1 \rightarrow B^P (f_0 a_0) (f_1 a_1)$$

$$(\Sigma A B)^P (a_0, b_0) (a_1, b_1) = \Sigma (a_2 : A^P a_0 a_1). B^P a_2 b_0 b_1$$

Informal Narya

$$\frac{A : \text{Type}}{A^P : A \rightarrow A \rightarrow \text{Type}} \qquad \frac{a : A}{a^P : A^P a a}$$

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$$(\lambda c. (a : A c) \rightarrow B(c, a))^P c_2 f_0 f_1 = (a_2 : A^P c_2 a_0 a_1) \rightarrow B^P (c_2, a_2) (f_0 a_0) (f_1 a_1)$$

Informal Narya

$$\frac{A : \text{Type}}{A^P : A \rightarrow A \rightarrow \text{Type}} \qquad \frac{a : A}{a^P : A^P a a}$$

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$$x^P = x_2 \qquad (\lambda x. t)^P = \lambda \{x_0\} \{x_1\} x_2. t^P \qquad (t u)^P = t^P u^P$$

Informal Narya

$$\frac{A : \text{Type}}{A^P : A \rightarrow A \rightarrow \text{Type}} \qquad \frac{a : A}{a^P : A^P a a}$$

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$$x^P = x_2 \qquad (\lambda x. t)^P = \lambda \{x_0\} \{x_1\} x_2. t^P \qquad (t u)^P = t^P u^P$$

$$\frac{a_{22} : A^{PP} a_{02} a_{12} a_{20} a_{21}}{\text{sym } a_{22} : A^{PP} a_{20} a_{21} a_{02} a_{12}} \qquad \text{sym (sym } a_{22}) = a_{22} \qquad \dots$$

Informal Narya

$$\frac{A : \text{Type}}{A^P : A \rightarrow A \rightarrow \text{Type}} \qquad \frac{a : A}{a^P : A^P a a}$$

$$\frac{R : A \rightarrow B \rightarrow \text{Type}}{\text{Gel } R : \text{Type}^P A B} \qquad \frac{A_2 : \text{Type}^P A_0 A_1}{A_2 : A_0 \rightarrow A_1 \rightarrow \text{Type}} \qquad \text{Gel } R a b \cong R a b$$

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$$x^P = x_2 \qquad (\lambda x. t)^P = \lambda \{x_0\} \{x_1\} x_2. t^P \qquad (t u)^P = t^P u^P$$

$$\frac{a_{22} : A^{PP} a_{02} a_{12} a_{20} a_{21}}{\text{sym } a_{22} : A^{PP} a_{20} a_{21} a_{02} a_{12}} \qquad \text{sym (sym } a_{22}) = a_{22} \qquad \dots$$

whiteboard example: polymorphic identity

Iteration \rightarrow induction

data \mathbb{N} : Type

zero : \mathbb{N}

suc : $\mathbb{N} \rightarrow \mathbb{N}$

ite $_{\mathbb{N}}$: $(A : \text{Type}) \rightarrow A \rightarrow (A \rightarrow A) \rightarrow \mathbb{N} \rightarrow A$

ite $_{\mathbb{N}}$ A z s zero = z

ite $_{\mathbb{N}}$ A z s (suc n) = s (ite $_{\mathbb{N}}$ A z s n)

Iteration \rightarrow induction

data \mathbb{N} : Type

zero : \mathbb{N}

suc : $\mathbb{N} \rightarrow \mathbb{N}$

data \mathbb{N}^P : Type^P \mathbb{N} \mathbb{N}

zero^P : \mathbb{N}^P zero zero

suc^P : \mathbb{N}^P n_0 n_1
 $\rightarrow \mathbb{N}^P$ (suc n_0) (suc n_1)

ite _{\mathbb{N}} : (A : Type) \rightarrow A \rightarrow (A \rightarrow A) \rightarrow $\mathbb{N} \rightarrow$ A

ite _{\mathbb{N}} A z s zero = z

ite _{\mathbb{N}} A z s (suc n) = s (ite _{\mathbb{N}} A z s n)

ite _{\mathbb{N}^P} : (A^P : Type^P \mathbb{N} \mathbb{N}) \rightarrow A^P zero zero

\rightarrow (A^P n_0 $n_1 \rightarrow$ A^P (suc n_0) (suc n_1)) \rightarrow \mathbb{N}^P n_0 n_1

\rightarrow A^P n_0 n_1

ite _{\mathbb{N}^P} A^P z^P s^P zero^P = z^P

ite _{\mathbb{N}^P} A^P z^P s^P (suc^P n^P) = s^P (ite _{\mathbb{N}^P} A^P z^P s^P n^P)

Iteration → induction

data \mathbb{N} : Type

zero : \mathbb{N}

suc : $\mathbb{N} \rightarrow \mathbb{N}$

data \mathbb{N}^P : Type^P \mathbb{N} \mathbb{N}

zero^P : \mathbb{N}^P zero zero

suc^P : \mathbb{N}^P n_0 n_1
 $\rightarrow \mathbb{N}^P$ (suc n_0) (suc n_1)

ite _{\mathbb{N}} : (A : Type) → A → (A → A) → \mathbb{N} → A

ite _{\mathbb{N}} A z s zero = z

ite _{\mathbb{N}} A z s (suc n) = s (ite _{\mathbb{N}} A z s n)

ite _{\mathbb{N}^P} : (A^P : Type^P \mathbb{N} \mathbb{N}) → A^P zero zero
 \rightarrow (A^P n_0 n_1 → A^P (suc n_0) (suc n_1)) → \mathbb{N}^P n_0 n_1
 \rightarrow A^P n_0 n_1

ite _{\mathbb{N}^P} A^P z^P s^P zero^P = z^P

ite _{\mathbb{N}^P} A^P z^P s^P (suc^P n^P) = s^P (ite _{\mathbb{N}^P} A^P z^P s^P n^P)

ind _{\mathbb{N}} : (P : $\mathbb{N} \rightarrow$ Type) → P zero → (P n → P (suc n)) → (n : \mathbb{N}) → P n

ind _{\mathbb{N}} P z^P s^P n :=

$$\begin{array}{c}
 \text{ite}_{\mathbb{N}^P} (\text{Glue } (n_0 \ n_1 \mapsto P \ n_0)) \quad \underbrace{z^P} \quad \underbrace{(n^P \mapsto s^P \ n^P)} \quad n^P \\
 \underbrace{\text{Glue } (n_0 \ n_1 \mapsto P \ n_0) \ \text{zero} \ \text{zero}}_{\cong P \ \text{zero}} \quad \underbrace{\text{Glue } (n_0 \ n_1 \mapsto P \ n_0) \ n_0 \ n_1}_{\cong P \ n_0} \quad \underbrace{\text{Glue } (n_0 \ n_1 \mapsto P \ n_0) \ (\text{suc } n_0) \ (\text{suc } n_1)}_{\cong P \ (\text{suc } n_0)} \\
 \hline
 : \text{Glue } (n_0 \ n_1 \mapsto P \ n_0) \ n \ n \cong P \ n
 \end{array}$$

Semi-cubical types*

codata $T : \text{Type}$

$\text{ty} : T \rightarrow \text{Type}$

$\text{rel} : (x : T) \rightarrow T^P x x$

codata $T^P : \text{Type}^P T T$

$\text{ty}^P : (T \rightarrow \text{Type})^P \text{ty ty}$

$\text{rel}^P : ((x : T) \rightarrow T^P x x)^P \text{rel rel}$

codata $T^{PP} : (\text{Type}^P T T)^P T^P T^P$

$\text{ty}^{PP} : ((T \rightarrow \text{Type})^P \text{ty ty})^P \text{ty}^P \text{ty}^P$

$\text{rel}^{PP} : (((x : T) \rightarrow T^P x x)^P \text{rel rel})^P \text{rel}^P \text{rel}^P$

...

Semi-cubical types*

codata $T : \text{Type}$

$\text{ty} : T \rightarrow \text{Type}$

$\text{rel} : (x : T) \rightarrow T^P x x$

codata $T^P : \text{Type}^P T T$

$\text{ty}^P : T^P x_0 x_1 \rightarrow \text{Type}^P (\text{ty } x_0) (\text{ty } x_1)$

$\text{rel}^P : (x_2 : T^P x_0 x_1) \rightarrow T^{PP} x_2 x_2 (\text{rel } x_0) (\text{rel } x_1)$

codata $T^{PP} : \text{Type}^{PP} T^P T^P T^P T^P$

$\text{ty}^{PP} : T^{PP} x_{02} x_{12} x_{20} x_{21} \rightarrow \text{Type}^{PP} (\text{ty}^P x_{02}) (\text{ty}^P x_{12}) (\text{ty}^P x_{20}) (\text{ty}^P x_{21})$

$\text{rel}^{PP} : (x_{22} : T^{PP} x_{02} x_{12} x_{20} x_{21}) \rightarrow T^{PPP} x_{22} x_{22} (\text{rel}^P x_{02}) (\text{rel}^P x_{12}) (\text{rel}^P x_{20}) (\text{rel}^P x_{21})$

...

Fibrancy

codata isFib : Type → Type

coe : isFib^P A₂ f₀ f₁ → A₀ → A₁

coh : (f₂ : isFib^P A₂ f₀ f₁)(a₀ : A₀) → A₂ a₀ (coe f₂ a₀)

coe⁻¹ : isFib^P A₂ f₀ f₁ → A₁ → A₀

coh⁻¹ : (f₂ : isFib^P A₂ f₀ f₁)(a₁ : A₁) → A₂ (coe⁻¹ f₂ a₁) a₁

id : (f₂ : isFib^P A₂ f₀ f₁)(a₀ : A₀)(a₁ : A₁) → isFib (A₂ a₀ a₁)

Fibrancy

codata isFib : Type → Type

coe : isFib^P A₂ f₀ f₁ → A₀ → A₁

coh : (f₂ : isFib^P A₂ f₀ f₁)(a₀ : A₀) → A₂ a₀ (coe f₂ a₀)

coe⁻¹ : isFib^P A₂ f₀ f₁ → A₁ → A₀

coh⁻¹ : (f₂ : isFib^P A₂ f₀ f₁)(a₁ : A₁) → A₂ (coe⁻¹ f₂ a₁) a₁

id : (f₂ : isFib^P A₂ f₀ f₁)(a₀ : A₀)(a₁ : A₁) → isFib (A₂ a₀ a₁)

Fib := Σ Type isFib

Fibrancy

codata isFib : Type → Type

coe : isFib^P A₂ f₀ f₁ → A₀ → A₁

coh : (f₂ : isFib^P A₂ f₀ f₁)(a₀ : A₀) → A₂ a₀ (coe f₂ a₀)

coe⁻¹ : isFib^P A₂ f₀ f₁ → A₁ → A₀

coh⁻¹ : (f₂ : isFib^P A₂ f₀ f₁)(a₁ : A₁) → A₂ (coe⁻¹ f₂ a₁) a₁

id : (f₂ : isFib^P A₂ f₀ f₁)(a₀ : A₀)(a₁ : A₁) → isFib (A₂ a₀ a₁)

Fib := Σ Type isFib

$A : \text{Type} \quad f : \text{isFib } A$

$f^P : \text{isFib}^P A^P f f$

$\text{coe } f^P : A \rightarrow A$

Two-dimensional Kan fillers

codata isFib : Type → Type

coe : isFib^P A₂ f₀ f₁ → A₀ → A₁ (also in the other direction)

coh : (f₂ : isFib^P A₂ f₀ f₁)(a₀ : A₀) → A₂ a₀ (coe f₂ a₀)

id : (f₂ : isFib^P A₂ f₀ f₁)(a₀ : A₀)(a₁ : A₁) → isFib (A₂ a₀ a₁)

Two-dimensional Kan fillers

codata isFib : Type → Type

coe : isFib^P A₂ f₀ f₁ → A₀ → A₁ (also in the other direction)

coh : (f₂ : isFib^P A₂ f₀ f₁)(a₀ : A₀) → A₂ a₀ (coe f₂ a₀)

id : (f₂ : isFib^P A₂ f₀ f₁)(a₀ : A₀)(a₁ : A₁) → isFib (A₂ a₀ a₁)

id^P : (f₂₂ : isFib^{PP} A₂₂ f₀₂ f₁₂ f₂₀ f₂₁)(a₀₂ : A₀₂ a₀₀ a₀₁)(a₁₂ : A₁₂ a₁₀ a₁₁)
→ isFib^P (A₂₂ a₀₂ a₁₂) (id f₂₀ a₀₀ a₁₀) (id f₂₁ a₀₁ a₁₁)

Two-dimensional Kan fillers

codata isFib : Type → Type

coe : isFib^P A₂ f₀ f₁ → A₀ → A₁ (also in the other direction)

coh : (f₂ : isFib^P A₂ f₀ f₁)(a₀ : A₀) → A₂ a₀ (coe f₂ a₀)

id : (f₂ : isFib^P A₂ f₀ f₁)(a₀ : A₀)(a₁ : A₁) → isFib (A₂ a₀ a₁)

id^P : (f₂₂ : isFib^{PP} A₂₂ f₀₂ f₁₂ f₂₀ f₂₁)(a₀₂ : A₀₂ a₀₀ a₀₁)(a₁₂ : A₁₂ a₁₀ a₁₁)
→ isFib^P (A₂₂ a₀₂ a₁₂) (id f₂₀ a₀₀ a₁₀) (id f₂₁ a₀₁ a₁₁)

$$\frac{\frac{f : \text{isFib } A}{f^{\text{PP}} : \text{isFib}^{\text{PP}} A^{\text{PP}} f^{\text{P}} f^{\text{P}} f^{\text{P}} f^{\text{P}}}}{f^{\text{PP}} : \text{isFib}^{\text{PP}} A^{\text{PP}} f^{\text{P}} f^{\text{P}} f^{\text{P}} f^{\text{P}}}}{\frac{a_{02} : A^{\text{P}} a_{00} a_{01} \quad a_{12} : A^{\text{P}} a_{10} a_{11}}{\text{coe} (\text{id}^{\text{P}} f^{\text{PP}} a_{02} a_{12}) : A^{\text{P}} a_{00} a_{10} \rightarrow A^{\text{P}} a_{01} a_{11}}}}$$

Two-dimensional Kan fillers

codata isFib : Type → Type

coe : isFib^P A₂ f₀ f₁ → A₀ → A₁ (also in the other direction)

coh : (f₂ : isFib^P A₂ f₀ f₁)(a₀ : A₀) → A₂ a₀ (coe f₂ a₀)

id : (f₂ : isFib^P A₂ f₀ f₁)(a₀ : A₀)(a₁ : A₁) → isFib (A₂ a₀ a₁)

id^P : (f₂₂ : isFib^{PP} A₂₂ f₀₂ f₁₂ f₂₀ f₂₁)(a₀₂ : A₀₂ a₀₀ a₀₁)(a₁₂ : A₁₂ a₁₀ a₁₁)
 → isFib^P (A₂₂ a₀₂ a₁₂) (id f₂₀ a₀₀ a₁₀) (id f₂₁ a₀₁ a₁₁)

$$\frac{\frac{f : \text{isFib } A}{f^{\text{PP}} : \text{isFib}^{\text{PP}} A^{\text{PP}} f^{\text{P}} f^{\text{P}} f^{\text{P}} f^{\text{P}}}}{f^{\text{PP}} : \text{isFib}^{\text{PP}} A^{\text{PP}} f^{\text{P}} f^{\text{P}} f^{\text{P}} f^{\text{P}}}}{\text{coe } (\text{id}^{\text{P}} f^{\text{PP}} a_{02} a_{12}) : A^{\text{P}} a_{00} a_{10} \rightarrow A^{\text{P}} a_{01} a_{11}}$$

$$\begin{array}{ccccc} & a_{01} & & a_{11} & \\ & \uparrow & & \uparrow & \\ a_{02} & & \Uparrow & & a_{12} \\ & a_{00} & \xrightarrow{\quad} & a_{10} & \\ & & a_{20} & & \end{array}$$

Identity type

$(A, f) : \text{Fib}$

$\text{Id}_{A,f} : A \rightarrow A \rightarrow \text{Fib}$

$\text{Id}_{A,f} a_0 a_1 := (A^{\text{P}} a_0 a_1, \text{id } f^{\text{P}} a_0 a_1)$

Identity type

$(A, f) : \text{Fib}$

$\text{Id}_{A,f} : A \rightarrow A \rightarrow \text{Fib}$

$\text{Id}_{A,f} a_0 a_1 := (A^P a_0 a_1, \text{id } f^P a_0 a_1)$

$B : A \rightarrow \text{Fib}$

$\text{snd} \circ B : (a : A) \rightarrow \text{isFib } (B a) \quad a_2 : A^P a_0 a_1$

$(\text{snd} \circ B)^P a_2 : \text{isFib}^P (B^P a^P) (\text{snd } (B a_0)) (\text{snd } (B a_1))$

$\text{transport}_B a_2 := \text{coe } ((\text{snd} \circ B)^P a_2) : \text{fst } (B a_0) \rightarrow \text{fst } (B a_1)$

Identity type

$(A, f) : \text{Fib}$

$\text{Id}_{A,f} : A \rightarrow A \rightarrow \text{Fib}$

$\text{Id}_{A,f} a_0 a_1 := (A^P a_0 a_1, \text{id } f^P a_0 a_1)$

$B : A \rightarrow \text{Fib}$

$\text{snd} \circ B : (a : A) \rightarrow \text{isFib } (B a) \quad a_2 : A^P a_0 a_1$

$(\text{snd} \circ B)^P a_2 : \text{isFib}^P (B^P a^P) (\text{snd } (B a_0)) (\text{snd } (B a_1))$

$\text{transport}_B a_2 := \text{coe } ((\text{snd} \circ B)^P a_2) : \text{fst } (B a_0) \rightarrow \text{fst } (B a_1)$

Assume $a : A$.

$(\Sigma(x : A).A^P a x)^P (a, a^P) (b, e) = \Sigma(x_2 : A^P a b).A^{PP} a^P x_2 a^P e$

Identity type

$$(A, f) : \text{Fib}$$

$$\text{Id}_{A,f} : A \rightarrow A \rightarrow \text{Fib}$$

$$\text{Id}_{A,f} a_0 a_1 := (A^P a_0 a_1, \text{id } f^P a_0 a_1)$$

$$B : A \rightarrow \text{Fib}$$

$$\text{snd} \circ B : (a : A) \rightarrow \text{isFib}(B a) \quad a_2 : A^P a_0 a_1$$

$$(\text{snd} \circ B)^P a_2 : \text{isFib}^P (B^P a^P) (\text{snd}(B a_0)) (\text{snd}(B a_1))$$

$$\text{transport}_B a_2 := \text{coe}((\text{snd} \circ B)^P a_2) : \text{fst}(B a_0) \rightarrow \text{fst}(B a_1)$$

Assume $a : A$.

$$(\Sigma(x : A). A^P a x)^P (a, a^P) (b, e) = \Sigma(x_2 : A^P a b). A^{PP} a^P x_2 a^P e$$

$$\begin{array}{ccc}
 a & \text{-----} & b \\
 a^P \uparrow & \uparrow & \uparrow e \\
 a & \xrightarrow{a^P} & a
 \end{array}$$

Fib closed under \top , \perp , \times , Σ , Π , Bool

F : Type \rightarrow Type

coe : $F^P C_2 f_0 f_1 \rightarrow C_0 \rightarrow C_1$ (also in other dir.)

coh : $(f_2 : F^P C_2 f_0 f_1)(c_0 : C_0) \rightarrow C_2 c_0 (coe f_2 c_0)$

id : $F^P C_2 f_0 f_1 \rightarrow (c_0 : C_0)(c_1 : C_1) \rightarrow F (C_2 c_0 c_1)$

$gen_{isFib} : (A : Type) \rightarrow F A \rightarrow isFib A$

Fib closed under \top , \perp , \times , Σ , Π , Bool

$$\begin{array}{l} F : \text{Type} \rightarrow \text{Type} \\ \text{coe} : F^{\text{P}} C_2 f_0 f_1 \rightarrow C_0 \rightarrow C_1 \quad (\text{also in other dir.}) \\ \text{coh} : (f_2 : F^{\text{P}} C_2 f_0 f_1)(c_0 : C_0) \rightarrow C_2 c_0 (\text{coe } f_2 c_0) \\ \text{id} : F^{\text{P}} C_2 f_0 f_1 \rightarrow (c_0 : C_0)(c_1 : C_1) \rightarrow F (C_2 c_0 c_1) \\ \hline \text{gen}_{\text{isFib}} : (A : \text{Type}) \rightarrow F A \rightarrow \text{isFib } A \end{array}$$

whiteboard

Fib closed under \top , \perp , \times , Σ , Π , Bool

$$\begin{array}{l} F : \text{Type} \rightarrow \text{Type} \\ \text{coe} : F^{\text{P}} C_2 f_0 f_1 \rightarrow C_0 \rightarrow C_1 \quad (\text{also in other dir.}) \\ \text{coh} : (f_2 : F^{\text{P}} C_2 f_0 f_1)(c_0 : C_0) \rightarrow C_2 c_0 (\text{coe } f_2 c_0) \\ \text{id} : F^{\text{P}} C_2 f_0 f_1 \rightarrow (c_0 : C_0)(c_1 : C_1) \rightarrow F (C_2 c_0 c_1) \\ \hline \text{gen}_{\text{isFib}} : (A : \text{Type}) \rightarrow F A \rightarrow \text{isFib } A \end{array}$$

whiteboard

Lemma: $\cong^{\text{P}} A_2 B_2 e_0 e_1 \rightarrow (a_0 : A_0)(a_1 : A_1) \rightarrow A_2 a_0 a_1 \cong B_2 (e_0 a_0) (e_1 a_1)$.

Fib closed under W and M

S : Type

P : $S \rightarrow \text{Type}$

W : Type

sup : $(s : S) \rightarrow (P s \rightarrow W) \rightarrow W$

Fib closed under W and M

S : Type

P : $S \rightarrow$ Type

W : Type

$\text{sup} : (s : S) \rightarrow (P s \rightarrow W) \rightarrow W$

Indexed W-types:

I : Type

S : Type

P : $S \rightarrow$ Type

j : $\{s : S\} \rightarrow P s \rightarrow I$

k : $S \rightarrow I$

W : $I \rightarrow$ Type

$\text{sup} : (s : S) \rightarrow ((p : P s) \rightarrow W (j p))$

$\rightarrow W (k s)$

Non-uniform parameterised W-types:

I : Type

S : $I \rightarrow$ Type

P : $\{i : I\} \rightarrow S i \rightarrow$ Type

k : $\{i : I\} \{s : S i\} \rightarrow P i s \rightarrow I$

W ($i : I$) : Type

$\text{sup} : (s : S i) \rightarrow ((p : P s) \rightarrow W (k p))$

$\rightarrow W i$

Univalence

$$\text{is11 } (R : A \rightarrow B \rightarrow \text{Fib}) := ((a : A) \rightarrow \text{isContr } (\Sigma(b : B).R a b)) \times \\ ((b : B) \rightarrow \text{isContr } (\Sigma(a : A).R a b))$$

Univalence

$$\text{is11 } (R : A \rightarrow B \rightarrow \text{Fib}) := ((a : A) \rightarrow \text{isContr } (\Sigma(b : B).R a b)) \times \\ ((b : B) \rightarrow \text{isContr } (\Sigma(a : A).R a b))$$

codata isBisim $(R : A \rightarrow B \rightarrow \text{Fib}) : \text{Type}$

coe : isBisim $R \rightarrow A \rightarrow B$ (also in the other direction)

coh : $(x : \text{isBisim } R)(a : A) \rightarrow R a (\text{coe } x a)$

id : $\text{isBisim}^R R_2 x_0 x_1 \rightarrow (r_0 : R_0 a_0 b_0)(r_1 : R_1 a_1 b_1) \rightarrow \text{isBisim } (a_2 b_2 \mapsto R_2 a_2 b_2 r_0 r_1)$

Univalence

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$\text{codata isBisim } (R : A \rightarrow B \rightarrow \text{Fib}) : \text{Type}$

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$$\text{is11} \longrightarrow \text{isBisim} \begin{array}{c} \xrightarrow{\text{Gel } R} \\ \xleftarrow{R := \text{id}} \end{array} \text{Fib}^P A_0 A_1$$

Univalence

$$\text{is11 } (R : A \rightarrow B \rightarrow \text{Fib}) := ((a : A) \rightarrow \text{isContr } (\Sigma(b : B).R a b)) \times ((b : B) \rightarrow \text{isContr } (\Sigma(a : A).R a b))$$

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$$\frac{\begin{array}{ccc} B_0 & \xrightarrow{B_2} & B_1 \\ R_0, \text{is11} \uparrow & R_2, \text{is11}^P \uparrow & R_1, \text{is11} \uparrow \\ A_0 & \xrightarrow{A_2} & A_1 \end{array}}{\text{is11 } (a_2 b_2 \mapsto R_2 a_2 b_2 r_0 r_1)}$$
$$\text{is11} \longrightarrow \text{isBisim} \begin{array}{c} \xleftarrow{\text{Gel } R} \\ \xrightarrow{R := \text{id}} \end{array} \text{Fib}^P A_0 A_1$$

Summary

- ▶ Narya: a proof assistant with internal parametricity
 - ▶ polymorphic identity
 - ▶ iterator \rightarrow induction, with definitional computation rules
 - ▶ symmetric inductive and coinductive types
 - ▶ semi-cubical types*
 - ▶ higher coinductive types, fibrancy
 - ▶ fibrant families come with transport
 - ▶ \top , \perp , \times , Σ , Π , Bool, W, M fibrant
 - ▶ definitional univalence
- ▶ Future:
 - ▶ isFib Fib,
 - ▶ normalisation
 - ▶ higher inductive types

Uniformity

